A Perspective on Photonic Multi-Protocol Label Switching



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Advanced Network Architecture Research

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- 1. MPλS (MPLS based on WDM Lightpaths)
  - Implementation Issues and Challenging Problems
- 2. OC-MPLS (MPLS based on Optical Code)
  - Optical Implementations and Challenging Issues
- 3. Perspectives on MPλS and OC-MPLS
  - Advantages and Disadvantages



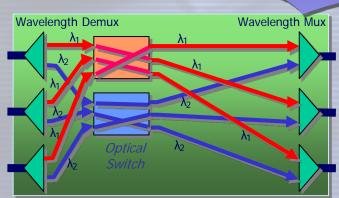
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# Mapping from Generic MPLS to Lambda MPLS (or GMPLS)

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LSR (Label Switching Router); Optical crossconnect directly connecting input wavelength to output wavelength



**Optical Crossconnect** 

LSP (Label-Switched Path); Wavelength path (Lightpath)

Ingress LSR; Maps from IP address to lambda

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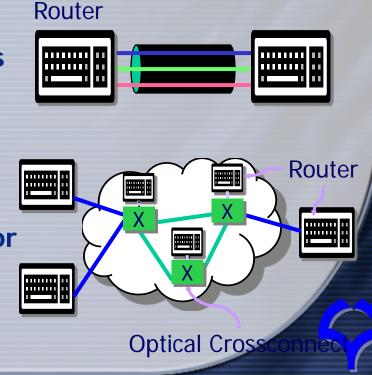
LDP (Label Distribution Protocol); Dimensioning by wavelength and routing assignment algorithm

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## **Photonic Internet Architecture**

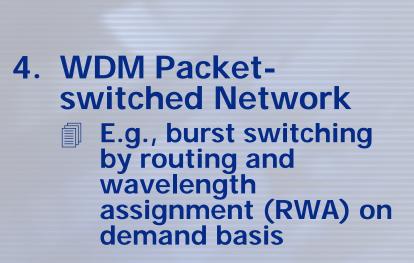
- Four Kinds of
   Architecture
  - 1. WDM link network
    - Connects adjacent routers by WDM (multiple wavelengths increase the bandwidth)
  - 2. WDM path network
  - Cut-through techniques for IP packets on established path provided by the underlying networks

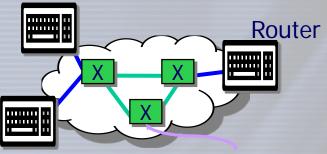




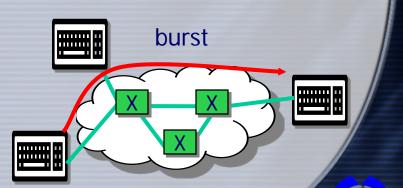
## Photonic Internet Architecture (Cont'd)

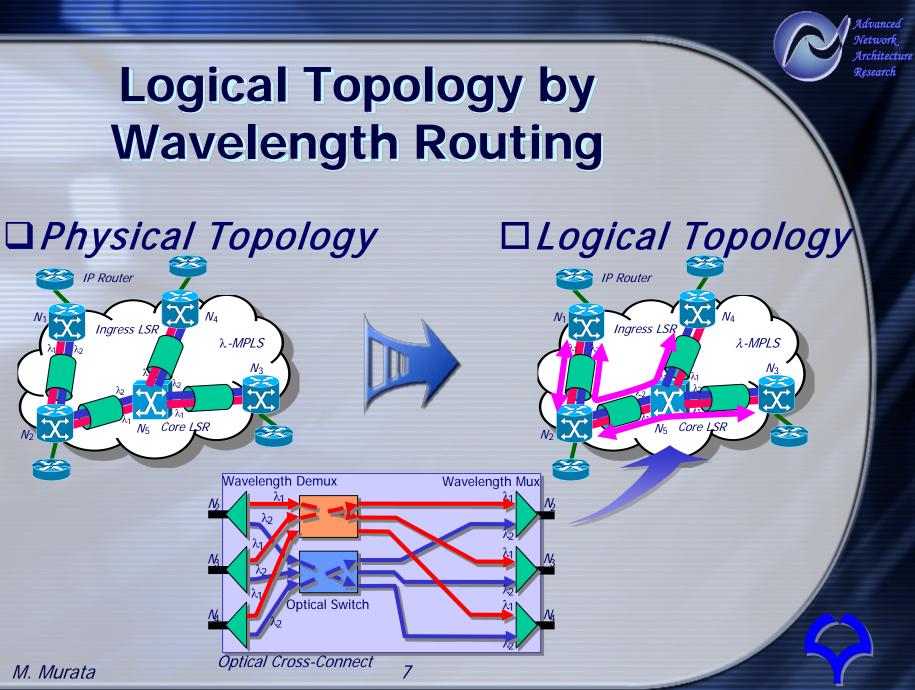
#### 3. WDM Path Network Lambda switching by MPLS technology (MPλS or GMPLS)





**Optical Cross-Connect** 

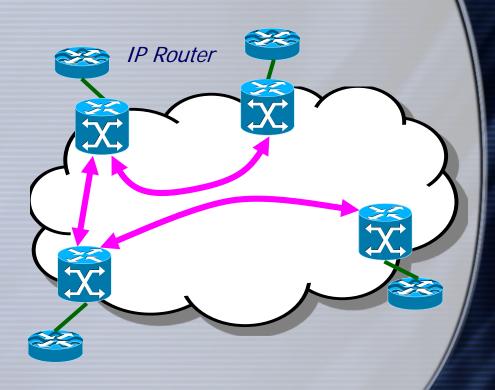






## **Logical View Provided to IP**

Redundant Network with Large Degrees Smaller number of hop-counts between end-nodes Decrease load for packet forwarding at the router Relief bottleneck at the router





## **Challenging Problems of MPλS**

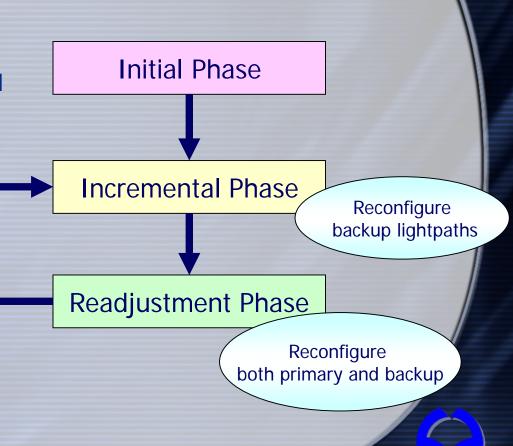
### Logical Topology Design Issues **O Past researches assume the traffic matrix is given** O Objective is maximization of wavelength utilization or minimization of the required # of wavelengths O Optimization problem is then solved by LP or heuristics Bottleneck at Ingress Nodes **O Bottleneck is shifted to the Ingress Nodes requiring** electronic processing Survivability **O IP and WDM Functional Partitioning or Integration**



## Incremental Capacity Dimensioning Approach

### Initial Phase

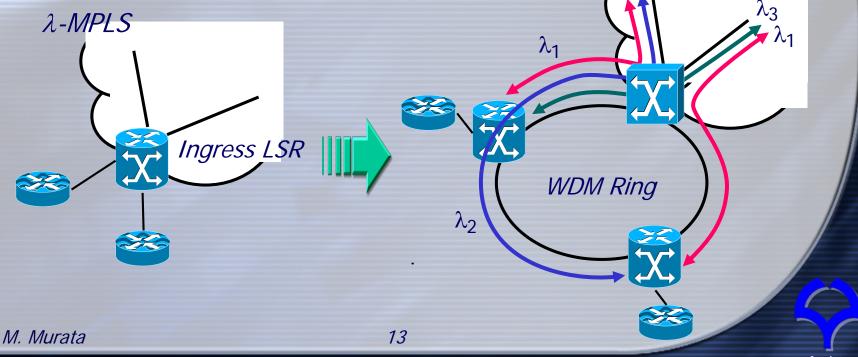
- Design Logical topology with given traffic demand
   Incremental Phase
  - primary lightpath is incrementally setup
  - O reconfigure backup lightpaths
- Readjustment Phase
  - All of the lightpath (including primary lightpath) is reconfigured





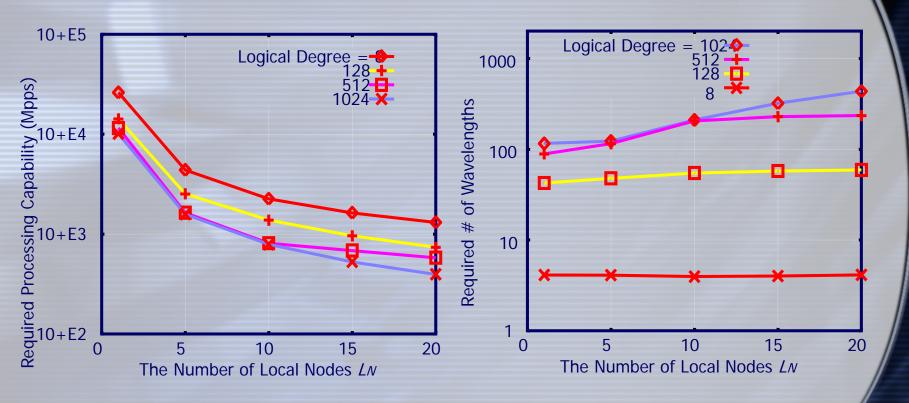
# Traffic Load Distribution by WDM Ring

### Distribute the traffic load by WDM ring at the Ingress node





### **Effects of Introducing WDM Ring**



Packet Processing Rate Required at Router

Required Number of Wave



## Do We Need More "Intelligent" WDM Network?

□ WDM network itself has network control capabilities

- **O** Routing function
  - IP also has it!
- O Congestion control function
  - TCP also has it!
  - TCP over ATM (ABR service class) is difficult to work well Parameter tuning of control parameters in ABR is not easy
- **O** Connection establishment
  - IP is connectionless
  - Multimedia application does not require 10Gbps channel
- **O Multi-layered Functionalities?**

Important is reliability



## **Functional Partitioning between IP and WDM?**

Reliability functionalities offered by two layers

 IP Layer: Routing
 WDM Layer: Path Protection and Restoration

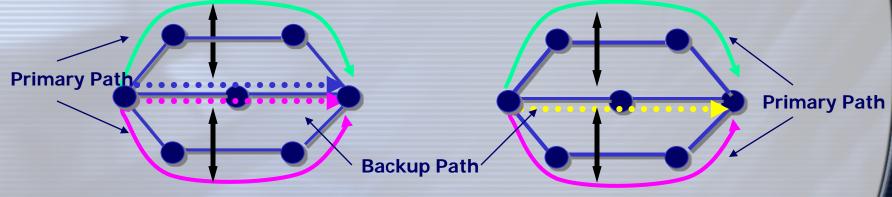
 WDM should provide its high-reliability mechanism to IP
 Protection mechanism

 link protection
 dedicated-path protection
 shared-path protection
 Network dimensioning is important to properly acquire the required capacity of IP paths (traffic glooming)
 Reconfiguration mechanism of logical topology by wavelength routing



### **WDM Protection**

Immediately switch to backup path on failure of nodes/links
 In the order of 10ms
 1:1 Protection vs. Many:1 Protection



#### **1:1 Protection**

**Many:1 Protection** 

Protection technique suitable to IP over WDM network?

- O IP has its own protection mechanism (i.e., routing) while it is slow
- **O** We want an effective usage of wavelengths
- O Many:1 protection is reasonable

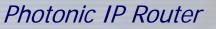
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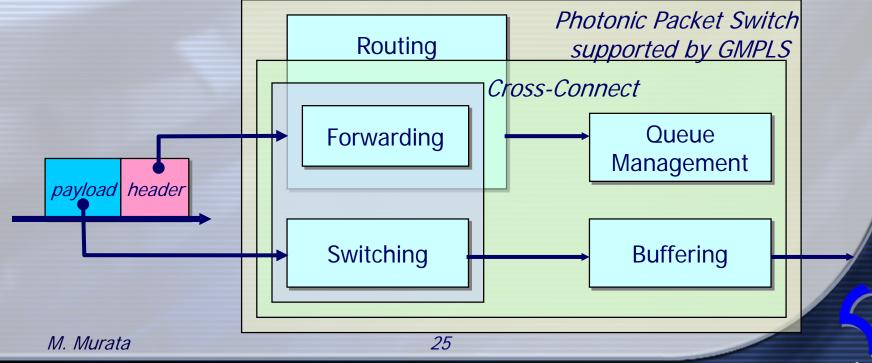


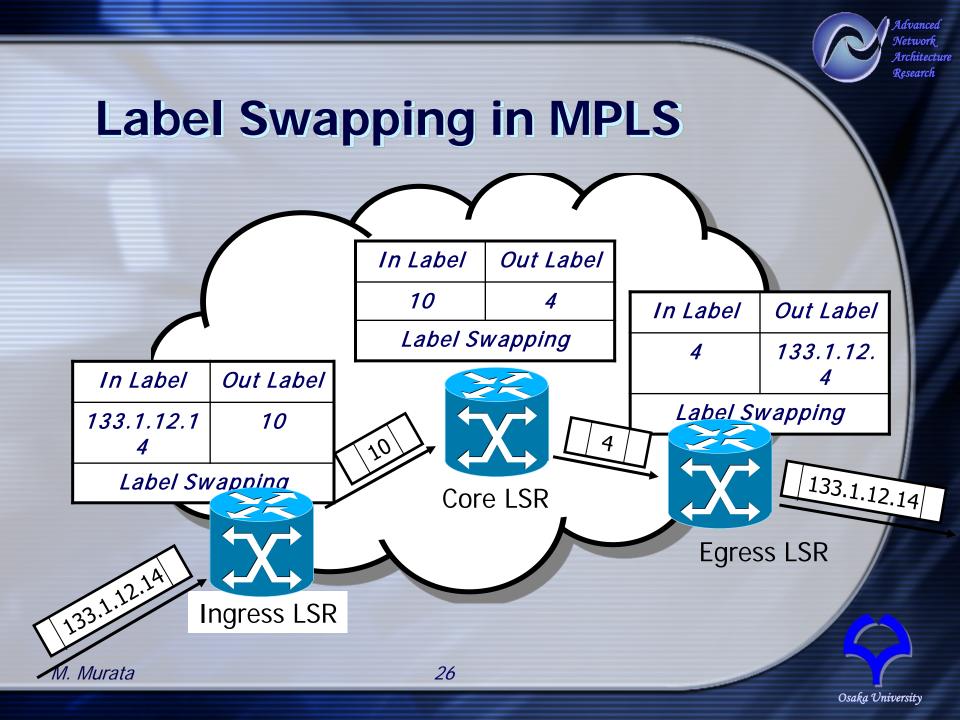
### **Cross-Connect, Switch and Router**

#### Photonic switch within MPLS requires

- O packet forwarding based on "label"
- O queue management based on "label"
- **O** packet switching and buffering





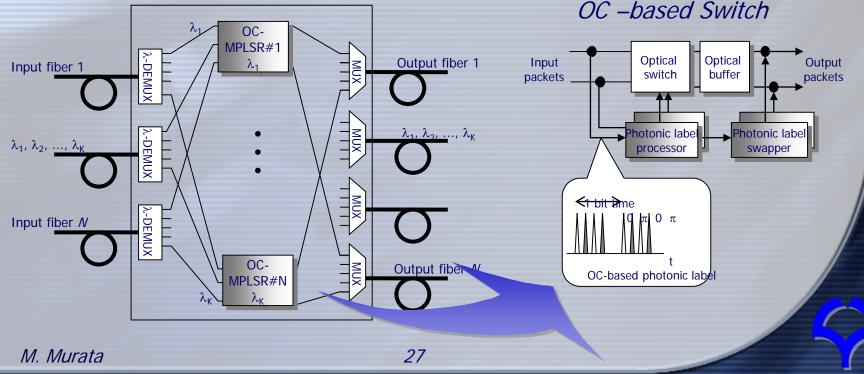




# **Optical Code based MPLS**

### Photonic label based on optical codes

#### OC-based Switching Node



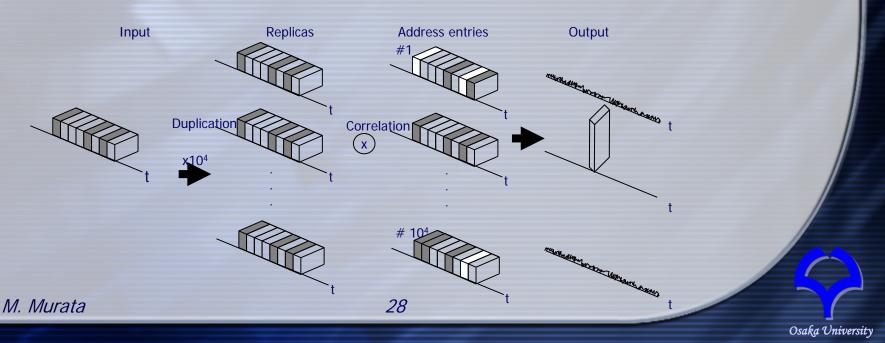


## **Photonic Label Processing**

#### Optical code by BPSK

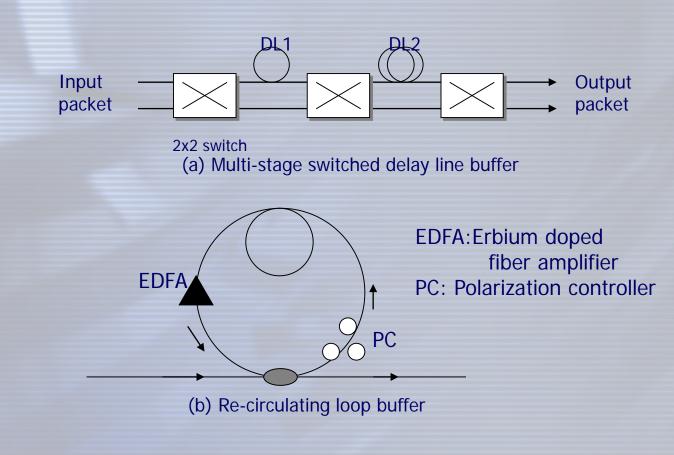
Photonic label processing in optical domain

- **O** photonic label is tapped from packet header
- **O** optically dupilcated by optical amplification
- **O** power-splitted as many copies as the count of label entries in the table
- optical correlations between the copies and the label entries is performed in parallel





### Implementations of Optical Buffers by Delay Lines



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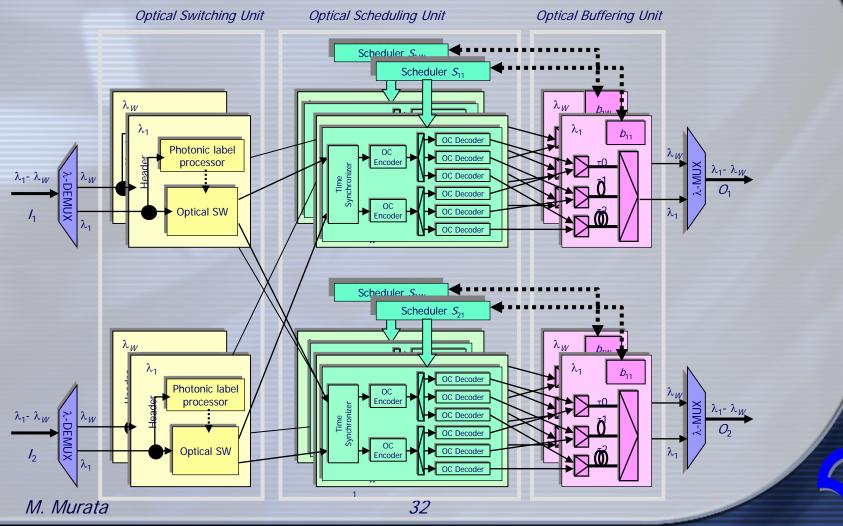


### Photonic Packet Switch for Asynchronously Arriving Variable Packets

 Buffer Scheduling **O Delayed Line Buffer** Variable-Length Packets O Counter for Buffer State;  $b_{ii}$  for output line *j* on wavelength **O** For each arrival of packet with length *x* b, ヤ+b, 端/D For each delay unit D  $b_{ii} + b_{ii} = 1$ □ For asynchronous arriving **O** Introduce packet sequencer



### Structure of OC-based Photonic Packet Switch Without Wavelength Conversion



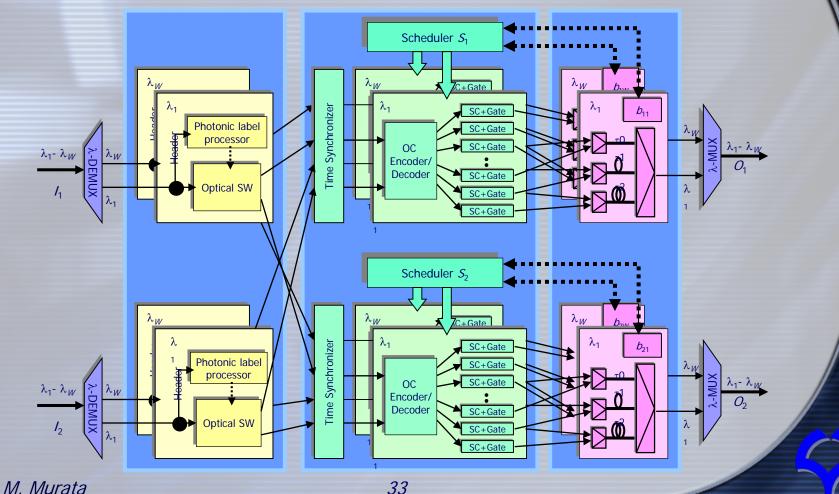


### Structure of OC-based Photonic Packet Switch With Wavelength Conversion

Optical Switching Unit

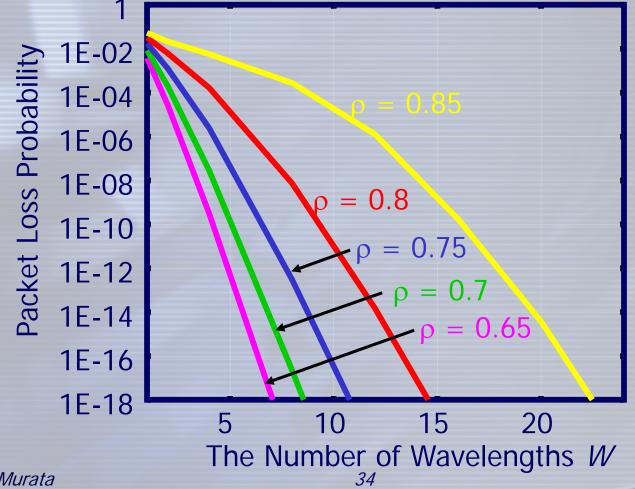
Optical Scheduling Unit

**Optical Buffering Unit** 





### **Performance of Proposed Switch**



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### **Problems of MPλS**

- The incremental capacity dimensioning is infeasible in the current logical topology design approach
  - O Network performance is heavily dependent on the logical topology design approach
- Unit of Path Granularity is Wavelength Capacity
  - O too large to accommodate the end-to-end traffic
  - O The capacity increase per each wavelength does not alleviate the problem
  - The increase in the number of wavelengths may help it. However, it requires the large-scaled of, e.g., 1,000x1,000 optical cross-connect
  - Flow aggregation at the Core LSR cannot be expected
    - The label exchange within the network poses the wavelength change at an optical node

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## **Advantages of OC-MPLS**

#### The granularity is "packet"

- **O** Allows a flexible network structure
- O Simplified packet switching can offer large capacity in the optical domain
- O An ATM-based MPLS protocol suite can be applied
- O Traffic engineering developed for MPLS is also utilized.
- OC-MPLS is capable of merging of packets by introducing optical buffering
  - O Attains an ultimate bandwidth efficiency
  - **O** MP $\lambda$ S is unable to realize it due to the coarse granularity.
  - The length of OC photonic label could be flexible
    - O The longer label could be used as the network layer header of the packet
    - Can be used both to assigned multiple flows from IP prefix to application-level flow
    - O Possible to offer QoS-enabled services
    - Optical codes are not only applicable to the exact match algorithm in the OC-MPLS but also applicable to the longest prefix match, and hence OCbased destination-based IP routing might be realizable.

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### Remaining Problems of Establishing OC-MPLS

- The switch fabric, constructed with photonic space switch and photonic buffer, has to be optimized to achieve the desired performance
  - Statistical multiplexing would work well by very huge bandwidth provided by OC techniques even if the packet buffer capacity is not large. However, a further research on the traffic engineering approach of MPLS is necessary under the conditions that the bandwidth is very large, but the packet buffer size is small.