

## Performance Analysis and Improvement of HighSpeed TCP with TailDrop/ARED Routers

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## Background

- Success of the Internet.
- Continuous and explosive growth of the Internet
  - Especially High-speed Long-delay network.
- New application – Data Grid, SANs, . . .
  - Hosts have gigabit-level network interface.
  - Perform backups, synchronize databases
- TCP (Transmission Control Protocol)
  - Traditional transport-layer protocol – TCP Reno used in the above applications
  - Can't achieve sufficient throughput.

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## Problems with TCP Reno<sup>[1]</sup>

- AIMD (Additive Increase Multiplicative Decrease) algorithm of TCP Reno in congestion avoidance phase:
  - No packet loss (AI): increase congestion window by one packet/RTT → too slowly
  - Packet loss (MD): decrease congestion window by half → too dramatically
- Example: For fully utilize a link of 10 Gbps with
  - Packet size: 1,500 bytes
  - RTT (Round Trip Time): 100 ms

This requires:

- Congestion window = 83,333 packets → Packet drop rate =  $2 \times 10^{-10}$

It is impossible with present optical fiber and router technology.

[1] S. Floyd, "HighSpeed TCP for large congestion windows," RFC 3649, December 2003.  
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## HighSpeed TCP – HSTCP<sup>[1]</sup>

- AIMD of HSTCP
  - AI: 1  $a(w)$
  - MD: 0.5  $b(w)$

[1] S. Floyd, "HighSpeed TCP for large congestion windows," RFC 3649, December 2003.  
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## Objectives

- Consider the mixture situation of TCP Reno and HSTCP
- Evaluate the performance of HighSpeed TCP
  - Simulation.
  - Metrics: Throughput and fairness against TCP Reno connections
- Propose 'Gentle' HighSpeed TCP
  - Higher throughput and better fairness than the original HighSpeed TCP

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## Evaluation

Case3: TCP Reno is used for  $S_1$  and HSTCP+SACK is used for  $S_2$ .

- TailDrop,  $\mu = 0.5$
- 3 Cases: Case1, Case2, Case3.

- Problems with HSTCP
  - Under-utilization (Case2)
  - Unfairness (Case3)
- Reasons
  - Buffer overflow leads that many packets are lost in bursty fashion
  - HSTCP still opens its CWND quickly even when packets are going to be stored in the router buffer

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## Gentle HighSpeed TCP – gHSTCP

- 2 modes in congestion avoidance phase
  - Observes the packet transmission time and its RTT:
    - If the positive correlation is recognized **Reno mode**
    - Otherwise **HSTCP Mode**

- gHSTCP:
  - Prevent bursty packet drop.
  - Provide suitable fairness to the competing TCP Reno flows.

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## Evaluation Case5: TCP Reno is used for S<sub>1</sub> and gHSTCP+SACK is used for S<sub>2</sub>.

- TailDrop, different access link bandwidth.
- 2 Cases: Case4, Case5.

- gHSTCP outperforms HSTCP in terms of:
  - Throughput.
  - Fairness.
- But the performance does not satisfactory, especially as the delay increases.

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## Why introduce ARED

- Problems with TailDrop
  - burst packet drop.
  - global synchronization.
  - Unfairness.
- RED (Random Early Detection) is recommended.
  - But control parameters in RED is highly sensitive to the network condition.
- ARED (Adaptive RED)
  - an improved version of RED.

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## ARED – Adaptive RED

- Difference between RED and ARED
  - Key point of ARED:
    - $max_p$  is dynamically adapted to keep the average queue size within the target queue boundaries according to network conditions.
- Target queue range:
 
$$[min_{th} + 0.4(max_{th} - min_{th}), min_{th} + 0.6(max_{th} - min_{th})]$$

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## Evaluation Case6: TCP Reno is used for S<sub>1</sub> and HSTCP is used for S<sub>2</sub>.

- ARED, different access link bandwidth.
- 2 Cases: Case6, Case7

- Even with ARED
  - Can't fully utilize link bandwidth
  - Fairness is not good.

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## Algorithm used by ARED

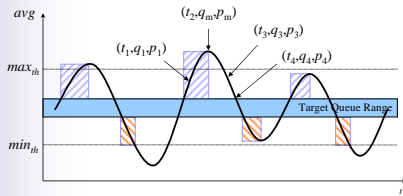
- Sketch of ARED
  - if  $(avg > target \text{ and } max_p < top)$  increase  $max_p$ :
  - if  $(avg < target \text{ and } max_p > bottom)$  decrease  $max_p$ :

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## Gentle Adaptive RED -- gARED

### Sketch of

if  $(avg > target \text{ and } avg > old\_avg \text{ and } max\_p < top)$   
 increase  $max\_p$ ;  
 if  $(avg < target \text{ and } avg < old\_avg \text{ and } min\_th < avg)$   
 decrease  $max\_p$ ;

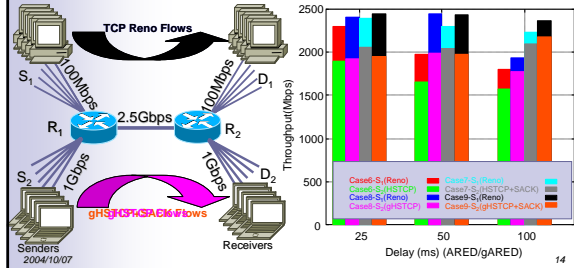


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## Evaluation

- Case9: TCP Reno is used for  $S_1$  and gHSTCP+SACK is used for  $S_2$ .
- gARED, different access link bandwidth.
- 2 Cases: Case8, Case9.



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## Conclusion

- gHSTCP+gARED can provide better performance in terms of fairness and throughput.
  - gHSTCP: two modes in congestion avoidance phase based on the changing trend of RTT
  - gARED: adapt  $max_p$  according to average queue length and its trend in variation
- Future works: further investigation, e.g.
  - Recover effectively from simultaneous packet losses
  - The impact on short-live traffic

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Thanks

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