Is tampered-TCP really effective for getting higher throughput in the Internet?

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- Currently, most Internet traffic is carried via TCP
  - The congestion control mechanism of TCP provides the fairness
  - Each OS implements this mechanism based on RFC
- There are TCP variants created by malicious users

We call these "*tampered-TCP*"

- Because TCP works at the end hosts, some users may modify its behavior

Considerable effects of tampered-TCP

1. Fairness among users may not be achieved
   - Tampered-TCP connections may selfishly monopolize the bandwidth

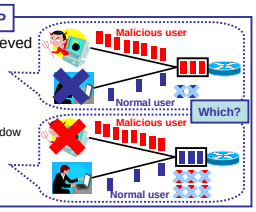
   ⇒ Normal TCPs suffer from low throughput
2. tampered-TCP may not work well

   ex. Augmenting increase ratio of congestion window

   ⇒ Congestion window becomes large rapidly

   ⇒ The number of loss packet increases and many timeouts occur

   ⇒ Tampered-TCP ***self-destructs***



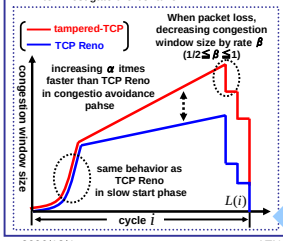
Our purpose

seeking the effects of tampered-TCP on a network

Tampered-TCP behavior and evaluation environment

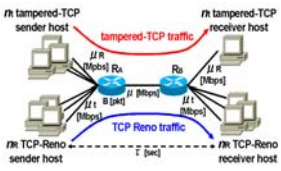
Behavior of tampered-TCP

- Considering a tampered-TCP which changes the increase and decrease ratio of congestion window size
  - Malicious users can expect the effect by easy modification
  - Researchers can use mathematical analysis to investigate the behavior



Network model and metric

Network model



Metric

$$Throughput\ ratio = \frac{(Throughput\ of\ \text{tampered-TCP})}{(Throughput\ of\ \text{TCP Reno})}$$

- When throughput ratio is greater than 1, tampered-TCP is said to work effectively.

Used for Mathematical analysis and simulation experiments

tampered-TCP evaluation (1)

Analysis

- Modeling cyclic change of window size triggered by a packet loss
- Considering a TCP property
  - More than 3 packets loss in a window

  ⇒ timeout occurs
- Considering a tampered-TCP property
  - The number of loss packets in a window ⇒ α

Timing of packet loss

$$\underbrace{n_R W_{Reno}(i, L(i))}_{\text{window size of TCP Reno}} + \underbrace{n_t W_t(i, L(i))}_{\text{window size of tampered-TCP}} > 2\tau\mu + B$$

Window size of tampered-TCP

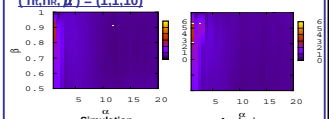
$$W_t(i,j) = \begin{cases} \beta^{\alpha} W_t(i-1, L(i-1)) + \alpha\, j & (\text{if } \alpha < 3) \\ \beta^{\alpha} W_t(i-1, L(i-1)) + \alpha \underbrace{[j - ss(i)]}_{\text{the number of RTT in slow start phase}} & (\text{if } \alpha \geq 3) \end{cases}$$

Throughput of tampered-TCP

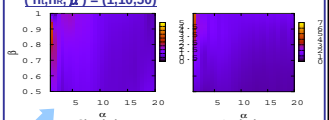
$$\rho_t = \frac{\sum_{i=1}^{\infty} \sum_{j=1}^{L(i)} W_t(i,j)}{\sum_{i=1}^{\infty} \sum_{j=1}^{L(i)} (\underbrace{Q(i,j)}_{\text{queuing delay}} + 2\tau)}$$

Throughput ratio change

(n_t,n_R, μ) = (1,1,10)



(n_t,n_R, μ) = (1,10,50)



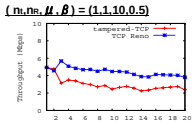
Analysis result is almost the same s simulation result

Tampered-TCP does not work for most of the (α, β)

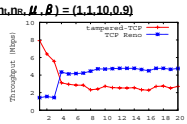
tampered-TCP evaluation (2)

Throughput change of each TCP connection

(n_t,n_R, μ , β) = (1,1,10,0.5)



(n_t,n_R, μ , β) = (1,1,10,0.9)

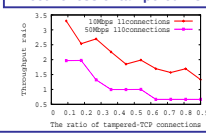


Sharp decrease when $\alpha \geqq 3$. More than 3 packets loss and timeouts occur.

Effective only when $\alpha < 3$. Sharp decrease when $\alpha \geqq 3$. More than 3 packets loss and timeouts occur.

Focus!

Effectiveness of tampered-TCP (α , β)=(2, 0.9)



The ratio of the number of tampered-TCP connections increases, the effectiveness sharply diminishes

tampered-TCP ***self-destructs***