#### Virtual Network Allocation for Fault Tolerance with Bandwidth Efficiency in a Multi-Tenant Data Center

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# Research background

- A data center (DC) for the laaS cloud computing
- serves virtual DC for multiple client organizations, i.e. tenants
   needs to host business-critical and mission-critical applications
- The virtual network (VN) for a tenant's virtual DC
- is an overlay network built by connecting VMs, based on VXLAN, etc
   has a topology independent of the physical substrate network (SN)
   should be appropriately assigned to the SN
  - to share the SN's resources effectively and tolerate SN failures

## • Goal:

ensuring high availability for the VN so that mission critical applications can be hosted on it

# **Research objectives**

- Mapping VNs to the shared physical SN is a kind of the *Virtual Network Embedding* problem
- Problems:

in a multi-tenant data center,

- nodes and links of VNs share a single component of the SN
- a failure of a single SN component can cause multiple simultaneous failures in a VN
  - significantly disrupts the services offered on the VN, as compared to a traditional network

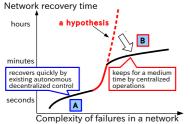
• Research objectives:

clarifying how the fault tolerance of a VN is affected by a SN failure, from the perspective of VN allocation

A hypothesis on the failure recovery time in a single VN

• A hypothesis: multiple simultaneous failures can lead to a longer recovery time in physical and virtual networks

• Proposal: switching from hot- to cold-standby recovery with reference to the failure complexity

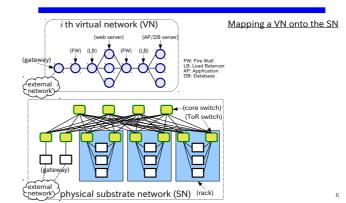


A Low complexity

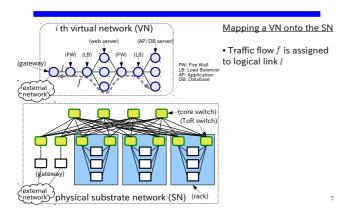
One or a few simultaneous failures
 VN recovers after a few seconds by
switching to hot-standby nodes and
links (VRRP, VMware FT, etc)

B High complexity • Many simultaneous failures > A centralized control force the failed nodes to be terminated and cold-standby nodes are alternatively booted (VMware HAT, etc)

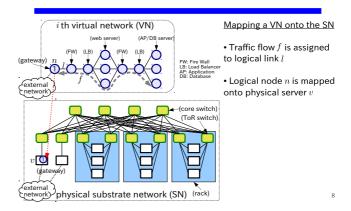
#### Network model for a multi-tenant data center



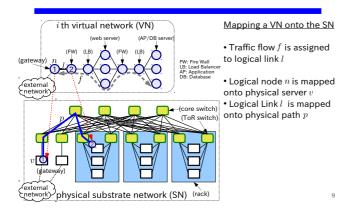
### Network model for a multi-tenant data center



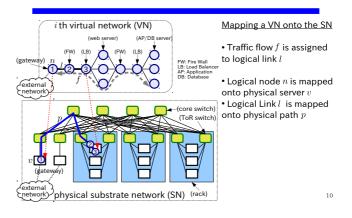
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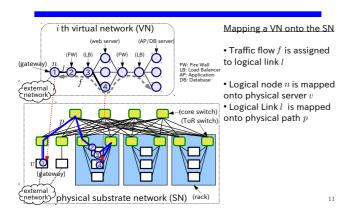
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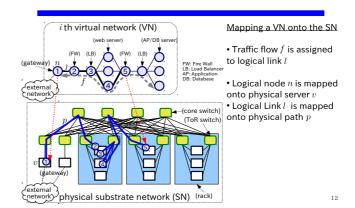
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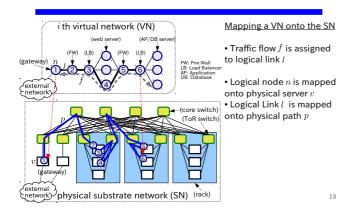
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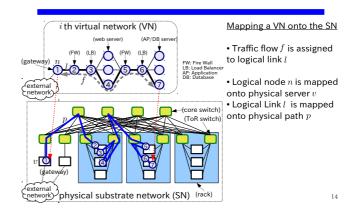
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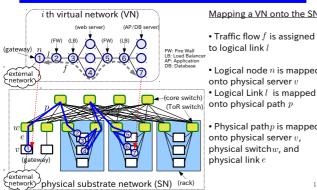
### Network model for a multi-tenant data center



### Network model for a multi-tenant data center



## Network model for a multi-tenant data center

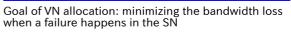


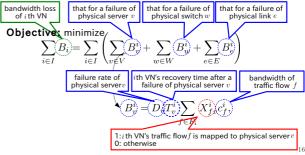
#### Mapping a VN onto the SN

- Traffic flow f is assigned
- Logical node n is mapped onto physical server v
- onto physical path  $\boldsymbol{p}$
- Physical path p is mapped onto physical server v, physical switchw, and

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# Objective





## Recovery time model of a single VN

- A failure of a single physical switch/link
- ▷ Recovery of the physical switch/link leads to recovery of the VN links
- Recovery time of the VN does not influenced by how the VN is embedded

 $2\mu$ 

- in the physical switch/link
- A failure of a single physical server  $T^{i}$ > VN should recover the VMs by utilizing
- its own failure-recovery mechanism Recovery time of the VN depends on
- how complicated the VN becomes
- failing simultaneously = the number of VMs
- assigned to the physical server

Subject to:

- rohibiting assigning more than θ VMs for ensuring the VMs' hot-standby recovery  $\sum_{n \in N_i} \mathbf{x}_{nv}^i \le \theta$
- *i* th VN's recovery time under physical server *v*'s failure cold standby hot



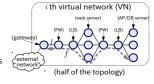
# Data center network for evaluation

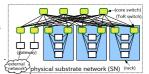
#### A single VN

- three-tier web serving architecture
   5.8 web and AP/DB servers,
- a total of 15.7 VMs on average > CPU cores per VM: 1
- average bandwidth demand
- from an external network:  $1.7 \times 10^8$  bit/s recovery time of a VM
   hot-standby: 4 s, cold-standby: 60 s

# The SN

- two-level fat-tree topology
   max configuration: 8 core switches,
- 16 ToR switches, and 120 physical servers > CPU cores/physical server: 32,
- bandwidth of each link: 1×10<sup>10</sup> bit/s available CPU cores: 3,360
- b failure rates physical server: 4/year, physical link/switch: 0.05/year (neglected)





# Overview of a single VN mapping

• VN embedding problem is NP-hard :

initially - Greedy Algorithm, refined - Tabu search • VN recovery time depends on  $\theta$  (threshold for switching hot- to cold-standby), which can not be defined in advance

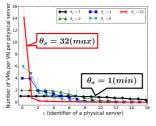
 $\triangleright \theta_s$  (a setting value of  $\theta$ ) is initially chosen  $\triangleright$  VN is allocated by using  $\theta_s$  , then evaluated for various values of  $\theta$ 

•  $\theta_s$  determines the *shape* of the VN  $\triangleright \; \theta_s = 1(min)$ 

- The VMs and logical links are scattered across many physical servers and links

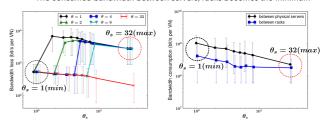
 $\triangleright \theta_s = 32(max)$ 

- All the VMs and links are consolidated in a few physical servers



### Trade-off between fault tolerance and physical bandwidth consumption

- $m{\cdot} heta_s = 1(min)$  : one VM to one physical server mapping The bandwidth loss is nearly the minimum for hot-standby recovery
   The consumed bandwidth between servers/racks reaches the maximum
- $\bullet \theta_s = 32(max)$  : many VMs to one physical server mapping  $\triangleright$  The bandwidth loss is nearly the maximum for cold-standby recovery The consumed bandwidth between servers/racks becomes the minimum

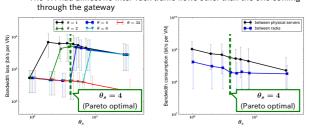


# Conclusion

• Minimizing the bandwidth loss of the VN while avoiding holding too many redundant core switches

VN Allocation Policy Derived from the Results

- Pareto optimality:  $\theta_s=4$   $\,^{\scriptscriptstyle b}$  Almost of the logical links were mapped onto the physical links
  - between the physical servers and ToR switches. > The VN had almost no inter-rack traffic flows other than the one coming



• The fault tolerance of each VN in an laaS data center ▷ Focusing on the situation of multiple simultaneous failures in each VN caused by a single physical failure > The trade-off between the bandwidth loss and the required bandwidth between physical servers

▷ Balancing by assigning every four VMs to a physical server, - the required bandwidth of the outside racks was minimized

• Future work

Investigation of resource allocation over WANs, i.e., in a hybrid cloud environment